

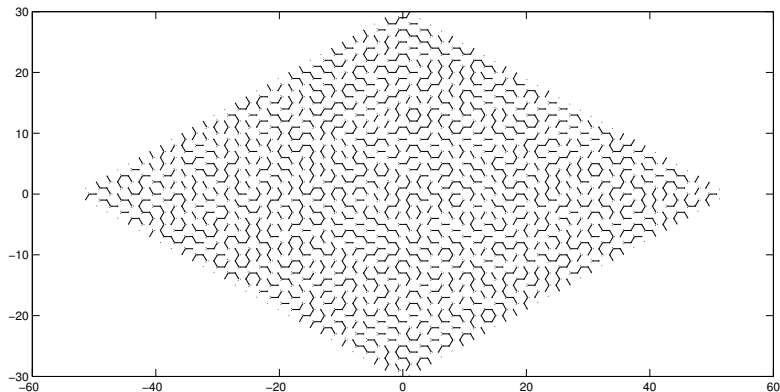
# 1-2 Model, Dimers and Clusters

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# 1-2 Model

- Probability measure on subgraphs  $\omega = (V, E_\omega)$  of hexagonal lattice  $\mathbb{H} = (V, E)$ , such that each vertex has 1 or 2 incident edges in  $\omega$



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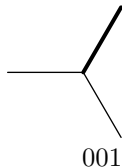
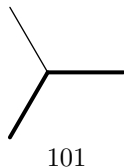
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- $Z$  is a normalizing constant called the partition function.

# Examples of local configurations



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- Each configuration and its complement have the same weight.
- The model is uniform if  $a = b = c = 1$ .



# History

- Computer scientists Schwartz and Bruck (2008) proposed the uniform 1-2 model (not-all-equal relation), as a graphical model whose partition function (total number of possible configurations) can be computed by computing determinants via holographic algorithm.

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- We introduced a generalized algorithm in a previous paper (2011), which could solve more models including the 1-2 model defined above.
- However, the holographic algorithm, although very general and beautiful, is not an efficient way to solve the 1-2 model.

# Correspondence: 1-2 model and dimers

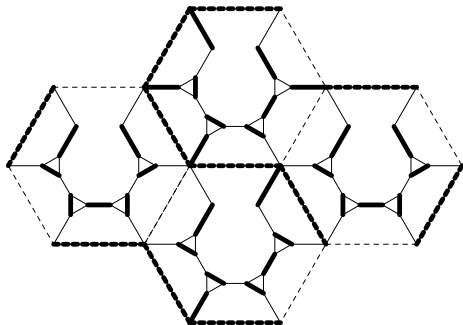


Figure: 1-2 Model and Dimers

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- $V \subset V_\Delta$ .
- For  $v \in V$ , incident edges of  $v$  in  $E_\Delta$  are bisectors of the angles of  $\mathbb{H}$  at  $v$ .
- On each face of  $\mathbb{H}$ , draw a small hexagon with vertices incident to the bisector edges.
- Remove the top most edge of each small hexagon.
- Change each degree-3 vertex of each small hexagon by a triangle.

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- A 1-2 model configuration and its complement correspond to the same perfect matching.
- Given appropriate edge weights to  $\mathbb{H}_\Delta$ , such a correspondence is measure-preserving.

# Clusters

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- An  $a$ -cluster is a connected set of vertices, all of which have  $a$ -configurations.
- In fact, for vertices in a single  $a$ -cluster, either all of them have the configuration 001, or all of them have the configuration 110, because the configuration 001 and 110 cannot appear on a pair of neighboring vertices.

# Existence of Phase Transition

## Theorem

*Fix  $b, c > 0$ , and use a large torus to approximate the infinite periodic graph. When  $a$  is sufficiently small, almost surely there is no infinite  $a$ -clusters; when  $a$  is large, the probability of the existence of infinite  $a$ -clusters is strictly positive, and the number of infinite  $a$ -clusters is at most one almost surely.*

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- Using a large deviation argument and results of determinantal processes, we prove that if  $V_0$  is an arbitrary set of vertices,

$$P(\text{All the vertices in } V_0 \text{ have } a\text{-configurations}) \leq (P_a)^{|V_0|}$$

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- We deduce that when  $a$  is small, almost surely there is no infinite  $a$ -clusters.

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- $S_p$ : the event that all the vertices in a box  $T$  centered at the origin have  $a$ -configurations, the number of boundary vertices of the box is no less than  $p$ .
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 & P(\text{an infinite } a \text{ - cluster appears at the origin}) \\
 & \geq P(\text{an infinite } a \text{ - cluster appears at the origin} | S_p) P(S_p) \\
 & = [1 - P(\text{no infinite } a \text{ - clusters at the origin} | S_p)] P(S_p) \\
 & \geq [1 - \sum_{q \geq p} \sum_{B_q} P(\text{none of vertices in } B_q \text{ have } a \text{ - configurations} | S_p)] \\
 & \quad \cdot P(S_p)
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## Sketch of Proof: 4



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- $\beta > 0$  is a constant independent of  $p$
- We deduce that when  $a$  is large, the probability that an infinite  $a$ -cluster appears at the origin is strictly positive.