

OPERADS IN ALGEBRAIC TOPOLOGY I

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INTRODUCTION

I'm really going to start from the beginning and build things up because operads are really important tools in algebraic topology.

Slogan. Operads encode n -ary operations and relations among them.

Motivating examples.

- (i) Associative monoids. Suppose I have a set X with a binary multiplication $\mu: X \times X \rightarrow X$ such that μ is associative:

$$\begin{array}{ccc}
 X \times X \times X & \xrightarrow{\mu \times X} & X \times X \\
 \downarrow X \times \mu & & \downarrow \mu \\
 X \times X & \xrightarrow{\mu} & X
 \end{array}$$

From the binary operation (X, μ) we get higher n -ary operations $X^{\times n} \rightarrow X$. In fact there exist $n!$ distinct n -ary operations, depending on how one permutes the inputs.

- (ii) Commutative monoids. Suppose we have (X, μ) as before where μ is associative and commutative: $\mu(x_1, x_2) = \mu(x_2, x_1)$. Again, this generates higher n -ary operations $X^{\times n} \rightarrow X$ but now there is a unique n -ary operation for each n .
- (iii) Based loop spaces. Let X be a based space and consider the space ΩX of based loops. We can define $\mu: \Omega X \times \Omega X \rightarrow \Omega X$ that sends a pair of loops to their concatenation: $(\lambda_1, \lambda_2) \mapsto \lambda_1 * \lambda_2$. This operation isn't strictly associative, but it is homotopy associative. There exists a based homotopy $H: \mu(\mu \times \Omega X) \simeq_* \mu(\Omega X \times \mu)$. Moreover, there exists a homotopy between the two homotopies from

$\mu(\mu(\mu \times \Omega X \times \Omega X))$ to $\mu(\Omega X \times \mu(\Omega X \times \Omega X \times \mu))$ induced by H .

$$\begin{array}{ccc}
 & (\lambda_1 * \lambda_2) * (\lambda_3 * \lambda_4) & \\
 \nearrow & & \searrow \\
 ((\lambda_1 * \lambda_2) * \lambda_3) * \lambda_4 & & \lambda_1 * (\lambda_2 * (\lambda_3 * \lambda_4)) \\
 \searrow & \Downarrow \text{htpy} & \nearrow \\
 (\lambda_1 * (\lambda_2 * \lambda_3)) * \lambda_4 & \longrightarrow & \lambda_1 * ((\lambda_2 * \lambda_3) * \lambda_4)
 \end{array}$$

- (iv) Double loop spaces: $\Omega^2 X$. This has two homotopy associative multiplications that satisfy an up-to-homotopy Eckmann-Hilton relation. This implies that they are homotopy commutative and that they are the same up to homotopy.

OPERADS: DEFINITIONS AND ELEMENTARY EXAMPLES

Let $(\mathcal{M}, \otimes, I)$ be a closed symmetric monoidal category. Here *closed* means that $- \otimes X: \mathcal{M} \rightarrow \mathcal{M}$ has a right adjoint $\text{Hom}(X, -): \mathcal{M} \rightarrow \mathcal{M}$. In particular, the counit of this adjunction defines a natural morphism

$$\text{ev}_X: \text{Hom}(X, Y) \otimes X \rightarrow Y$$

for all $X, Y \in \mathcal{M}$. I also want \mathcal{M} to be cocomplete.

Example. $(\mathbf{sSet}, \times, \{*\})$, $(\mathbf{Top}^{\text{nice}}, \times, \{*\})$, $(\mathbf{Ch}_R, \otimes, R)$.

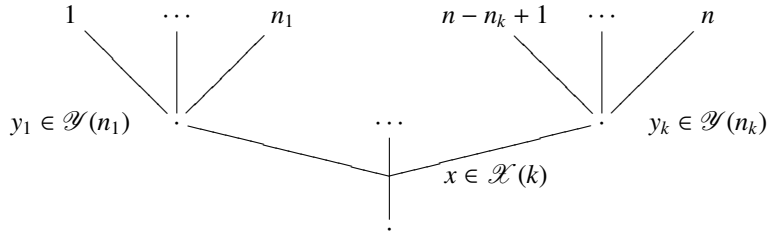
Definition. Let \mathcal{M}^{Σ} denote the category of *symmetric sequences* in \mathcal{M} . Objects are $\mathcal{X} = (\mathcal{X}(n))_{n \geq 0}$ where $\mathcal{X}(n) \in \mathcal{M}$ is an object that admits a Σ_n -action (i.e., there exists a homomorphism $\Sigma_n \rightarrow \text{Aut}(\mathcal{X}(n))$). A morphism $f: \mathcal{X} \rightarrow \mathcal{Y}$ is a collection of Σ_n -equivariant maps $f_n: \mathcal{X}(n) \rightarrow \mathcal{Y}(n)$ for all $n \geq 0$.

Remark. \mathcal{M}^{Σ} admits several monoidal structures. For example,

- the *level monoidal structure* \otimes with $(\mathcal{X} \otimes \mathcal{Y})(n) = \mathcal{X}(n) \otimes \mathcal{Y}(n)$ given the diagonal Σ_n -action. The unit is $\mathcal{C} = (\mathcal{C}(n))$, where $\mathcal{C}(n) = I$.
- the *graded monoidal structure*, the *matrix monoidal structure*
- the *composition monoidal structure* \circ , which is non-symmetric. For $n \geq 1$, define

$$(\mathcal{X} \circ \mathcal{Y})(n) = \coprod_{k \geq 1, \vec{n}=(n_1, \dots, n_k), \sum n_i=n} \mathcal{X}(k) \otimes_{\Sigma_k} (\mathcal{Y}(n_1) \otimes \dots \otimes \mathcal{Y}(n_k)) \otimes_{\Sigma_{n_1} \times \dots \times \Sigma_{n_k}} I[\Sigma_n]$$

where $I[\Sigma_n]$ is the tensor of $n!$ copies of the monoidal unit I . The unit for \circ is the symmetric sequence \mathcal{I} that has I in arity one and the initial object everywhere else. Here is a schematic picture:



Recall. A *monoid* in a monoidal category $(\mathcal{M}, \otimes, I)$ is $A \in \mathcal{M}$, a multiplication $\mu: A \otimes A \rightarrow A$, and a unit $\eta: I \rightarrow A$ so that

$$\begin{array}{ccc} A \otimes A \otimes A & \xrightarrow{\mu \otimes A} & A \otimes A \\ A \otimes \mu \downarrow & & \downarrow \mu \\ A \otimes A & \xrightarrow{\mu} & A \end{array} \qquad \begin{array}{ccc} A \cong A \otimes I & \xrightarrow{A \otimes \eta} & A \otimes A \xleftarrow{\eta \otimes A} I \otimes A \cong A \\ & \searrow = & \downarrow \mu \swarrow = \\ & & A \end{array}$$

Definition. An *operad* in \mathcal{M} is a monoid in \mathcal{M}^Σ with respect to the composition monoidal structure. I.e., (\mathcal{P}, μ, η) where $\mathcal{P} = (\mathcal{P}(n))_{n \geq 0}$, $\mu: \mathcal{P} \circ \mathcal{P} \rightarrow \mathcal{P}$, $\eta: I \rightarrow \mathcal{P}$. Here

$$\mu \leftrightarrow \{\mu_{k, \vec{n}}: \mathcal{P}(k) \otimes (\mathcal{P}(n_1) \otimes \cdots \otimes \mathcal{P}(n_k)) \rightarrow \mathcal{P}(n)\}_n$$

with appropriate equivariance conditions.

Example.

- (i) Let $X \in \mathcal{M}$. The *endomorphism operad* on X is $\text{End}(X)$ where $\text{End}(X)(n) = \text{Hom}(X^{\otimes n}, X)$ and $\mu: \text{End}(X) \circ \text{End}(X) \rightarrow \text{End}(X)$ corresponds to

$$\mu_{k, \vec{n}}: \text{Hom}(X^{\otimes k}, X) \otimes (\text{Hom}(X^{\otimes n_1}, X) \otimes \cdots \otimes \text{Hom}(X^{\otimes n_k}, X)) \rightarrow \text{Hom}(X^{\otimes n}, X)$$

given by $f \otimes (g_1 \otimes \cdots \otimes g_k) \mapsto f \circ (g_1 \otimes \cdots \otimes g_k)$. (This can also be defined abstractly using the adjunctions.)

- (ii) The *associative operad* $\text{As}(n) = I[\Sigma_n]$. Here

$$\Sigma_k \times (\Sigma_{n_1} \times \cdots \times \Sigma_{n_k}) \xrightarrow{\mu} \Sigma_n$$

is defined so that $\sigma(\tau, \tau_1, \dots, \tau_k)$ is the permutation of n letters partitioned into k boxes where each τ_i permutes the elements of a box and τ permutes the boxes.

- (iii) The *commutative operad* is $\text{Com} = \mathcal{C}$ with $\mathcal{C}(n) = I$. Here μ consists of the unit isomorphisms $I \otimes (I \otimes \cdots \otimes I) \xrightarrow{\cong} I$. This was the monoidal unit for the level monoidal structure.

ALGEBRAS

Definition. Let (\mathcal{P}, μ, η) be an operad in \mathcal{M} . A \mathcal{P} -*algebra* is an object $X \in \mathcal{M}$ together with an operad map $\phi: \mathcal{P} \rightarrow \text{End}(X)$.

You can think of this as being a representation of \mathcal{P} on X . What does this mean? The components of the map are Σ_n -equivariant maps

$$\phi_n: \mathcal{P}(n) \rightarrow \text{End}(X)(n) = \text{Hom}(X^{\otimes n}, X).$$

These transpose to maps

$$\phi_n^b: \mathcal{P}(n) \otimes_{\Sigma_n} X^{\otimes n} \rightarrow X.$$

Since ϕ is an operad map, there exist an ‘‘associativity’’ relation:

$$\begin{array}{ccc} & \mathcal{P}(k) \otimes (\mathcal{P}(n_1) \otimes \cdots \otimes \mathcal{P}(n_k)) \otimes X^{\otimes n} & \\ \mu \otimes X^{\otimes n} \swarrow & & \searrow \mathcal{P}(k) \otimes \phi_{n_1}^b \otimes \cdots \otimes \phi_{n_k}^b \\ \mathcal{P}(n) \otimes X^{\otimes n} & & \mathcal{P}(k) \otimes X^{\otimes k} \\ \phi_n^b \searrow & & \swarrow \phi_k^b \\ & X & \end{array}$$

Slogan. $\mathcal{P}(n)$ parametrizes the n -ary operations on the \mathcal{P} -algebra X .

Example.

(i) An As-algebra is a monoid

$$\text{As}(n) \otimes_{\Sigma_n} X^{\otimes n} \cong X^{\otimes n} \rightarrow X.$$

(ii) A Com-algebra is a commutative monoid

$$\text{Com}(n) \otimes_{\Sigma_n} X^{\otimes n} \cong X^{\otimes n} / \Sigma_n \rightarrow X.$$

Remark. A morphism of operads $\phi: \mathcal{P} \rightarrow \mathcal{Q}$ induces a functor $\phi^*: \mathbf{Alg}_{\mathcal{Q}} \rightarrow \mathbf{Alg}_{\mathcal{P}}$ by pullback of structure: $\mathcal{P} \rightarrow \mathcal{Q} \rightarrow \text{End}(X)$.

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